# Introduction to random hyperbolic Graphs

Dieter Mitsche

Colloquium, Univ. Potsdam, June 17, 2020



Part I: Motivation and model specification

### Random hyperbolic graphs (RHGs): Introduction

▶ Introduced by Krioukov, Papadopoulos, Kitsak, Vahdat, Boguñá [Phys. Rev. '10]

 Appeal: Replicate characteristic properties observed in "real world networks" or "complex networks"

Example of networks: Power grid

Internet

Social networks

Biological interaction networks

...

Typical properties: Sparse

Heterogeneous

Locally dense (exhibit clustering phenomena)

Small world Navigable

Scale-free (with exponent between 2 and 3)

...

### Random hyperbolic graphs (RHGs): Introduction

▶ Introduced by Krioukov, Papadopoulos, Kitsak, Vahdat, Boguñá [Phys. Rev. '10]

 Appeal: Replicate characteristic properties observed in "real world networks" or "complex networks"

Example of networks: Power grid

Internet

Social networks

Biological interaction networks

...

Typical properties: Sparse

Heterogeneous

Locally dense (exhibit clustering phenomena)

Small world Navigable

Scale-free (with exponent between 2 and 3)

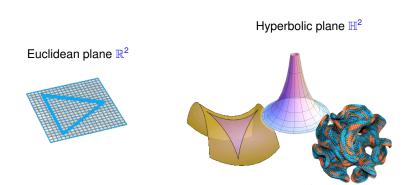
...

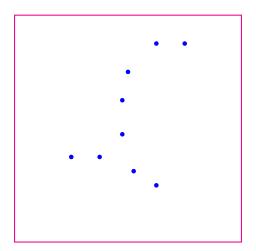
#### Susceptible to mathematical analysis!

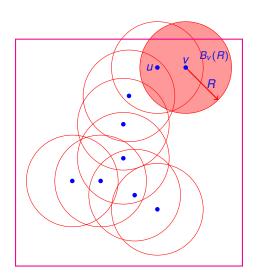


#### Informal definition of RHGs model

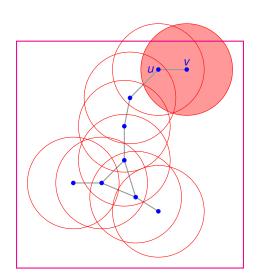
Like random geometric graphs but where the underlying space instead of being Euclidean is Hyperbolic.

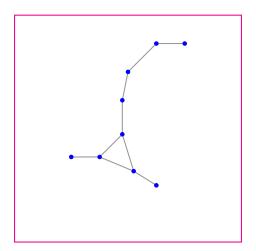




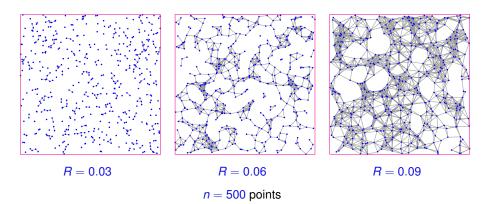






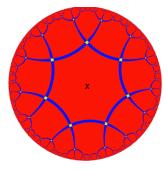


# Examples of random geometric graphs





#### Poincaré disk model of H<sup>2</sup>



[Rendered with KaleidoTile by J. Weeks]

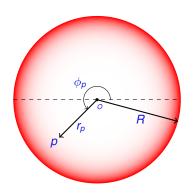
- ▶ ℍ² is represented as an open unit disk D
- Blue curves are geodesics (arcs of circles perpendicularly incident to D)
- Each heptagon has the same area
- Points in ∂D are at infinite distance from X
- Points at (Euclidean) distance y from X are at hyperbolic distance r from X where

$$r = \log \frac{1+y}{1-y}.$$

Space expands at exponential rate!

Continuous analogue of regular trees.

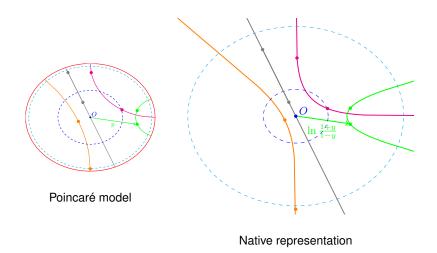
# Native representation of $\mathbb{H}^2$



 $B_O(R)$ : Ball of radius R centered at origin O with perimeter  $2\pi \sinh R = \Theta(e^R)$ .

- ▶  $\mathbb{H}^2$  is represented as  $\mathbb{R}^2$
- ► A point *p* is represented in polar coordinates
- ► r<sub>p</sub> is the hyperbolic distance between p and O

# Poincaré vs Native representation of $\mathbb{H}^2$



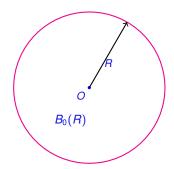
# Formal definition of RHG model: $G_{\alpha,\nu}(n)$

(Gugelmann, Panagiotou, Peter [ICALP'12])

#### Model parameters:

$$\alpha, \nu \in \mathbb{R}_+, \mathbf{n} \in \mathbb{N}_+.$$

Set 
$$R := 2 \log \frac{n}{\nu}$$
.



Choose an *n*-node graph G = (V, E) as follows:

- Each v ∈ V uniformly and independently in B<sub>O</sub>(R) according to some distribution depending on α
- ▶  $uv \in E$  iff  $u \in B_v(R)$ .

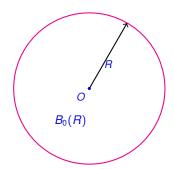
# Formal definition of RHG model: $G_{\alpha,\nu}(n)$

(Gugelmann, Panagiotou, Peter [ICALP'12])

#### Model parameters:

$$\alpha, \nu \in \mathbb{R}_+, n \in \mathbb{N}_+$$

Set  $R := 2 \log \frac{n}{\nu}$ .



#### Choose an *n*-node graph G = (V, E) as follows:

▶ Each  $v \in V$  chooses  $\phi_v \sim \text{Unif}[0, 2\pi)$  independent of  $r_v$  with density:

$$f(r) := \frac{\alpha}{C_{\alpha,R}} \sinh(\alpha r) \approx \alpha e^{-\alpha(R-r)}$$
 if  $0 \le r < R$  and 0 otherwise.

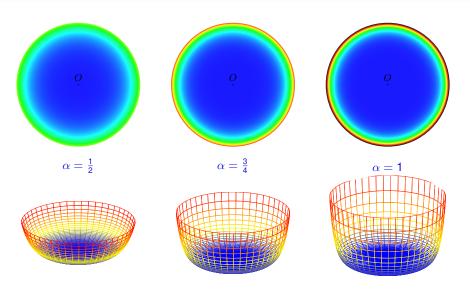
(Here,  $C_{\alpha,R}$  is a normalizing constant).

▶  $uv \in E$  iff  $u \in B_v(R)$ .



# Pdf of $(r_v, \phi_v)$ and its heat plot

(Colder colors correspond to smaller density)



# Calculating distances

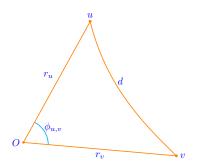
Hyperbolic distance from v to origin O, ... easy! Just  $r_v$ .

### Calculating distances

Hyperbolic distance from v to origin O, ... easy! Just  $r_v$ .

In general, use hyperbolic law of cosines

$$\cosh(d) = \cosh(r_u) \cosh(r_v) - \sinh(r_u) \sinh(r_v) \cos(\phi_{u,v}).$$



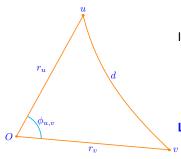


# Calculating distances

Hyperbolic distance from v to origin O, ... easy! Just  $r_v$ .

In general, use hyperbolic law of cosines

$$\cosh(d) = \cosh(r_u) \cosh(r_v) - \sinh(r_u) \sinh(r_v) \cos(\phi_{u,v}).$$



If 
$$d = R$$
 and  $r_u + r_v > R$ , then  $[GPP^*12]$ 

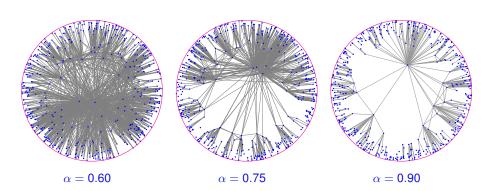
$$A_T(r, r) := 2a^{\frac{1}{2}(R-r_u-r_v)}(1 + R)(a^{R-r_u-r_v})$$

$$\begin{aligned} \theta_R(r_u, r_v) &:= 2e^{\frac{1}{2}(R - r_u - r_v)} (1 + \Theta(e^{R - r_u - r_v})) \\ &= \Theta(e^{\frac{1}{2}(R - r_u - r_v)}). \end{aligned}$$

**Lemma:** 
$$\phi_{u,v} \leq \theta_R(r_u, r_v) \iff d_{\mathbb{H}^2}(u, v) \leq R$$
.

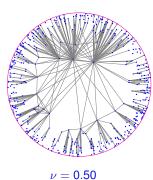
# Examples of RHGs

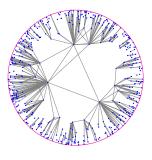
 $(\nu = 1 \text{ fixed}, n = 500)$ 



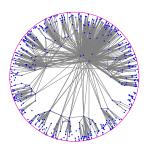
# Examples of RHGs

 $(\alpha = \frac{3}{4} \text{ fixed}, n = 500)$ 





 $\nu = 0.75$ 



 $\nu = 1.00$ 

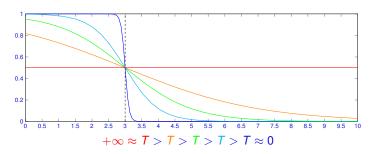


#### Soft version

Incorporates a temperature T and a probability of connecting u and v:

$$p(d) := \frac{1}{1 + e^{\frac{1}{2T}(d-R)}}$$

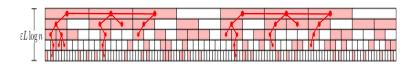
where  $d := d_{\mathbb{H}^2}(u, v)$  is the (hyperbolic) distance between  $u, v \in \mathbb{H}^2$ .



$$R = 3.0.$$



# Alternative representation of $\mathbb{H}^2$ using the halfplane model



Halfplane representation (FM18) on  $(-n, n) \times [0, R]$ 

- $ightharpoonup \phi: (r,\theta) 
  ightarrow (rac{e^{R/2}}{2},R-r)$
- $(x,y),(x',y') \in E \text{ iff } |x-x'| \leq \frac{e^{y+y'}}{2}$
- Advantage: easy to see that it converges locally to the infinite model

◆ロト 4個 ト 4 重 ト 4 重 ト 9 年 9 9 ○

# Nice, but who cares?

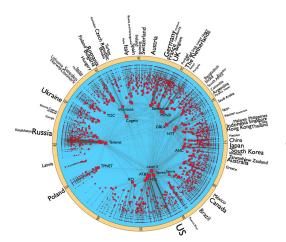
First model that "naturally" exhibits:

- ► Scale freeness, AND
- Non-negligible clustering.

But, what really drew attention ...

### Mapping of Internet's Autonomous Systems (ASs)

(2009 data collected by infrastructure developed by CAIDA)



[From Boguñá, Papadopoulus, Krioukov (Nat. Comm. '10)]

#### Data set:

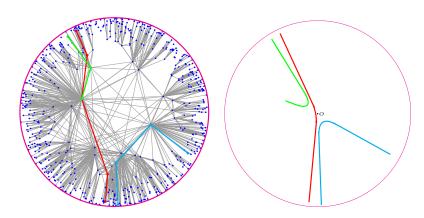
- 23,752 ASs
- ▶ 58,416 links
- Average degree 4.92

#### "Maximum Likelihood" fit:

- ▶  $\alpha = 0.55$
- ► R = 27
- ► Temperature T = 0.69



## **Greedy Forwarding**

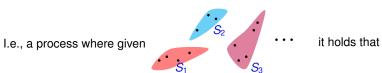


Papadopoulos et al. [INFOCOM 2010] in an experimental study (but without "real" data) report excellent stretch (average  $\sim$  1, max  $\sim$  1.4) and success ratio (0.99920 for  $\alpha \sim \frac{1}{2}$  to 0.92 for  $\alpha \sim 1$ , with  $\alpha, \nu$  as in the Internet).

Part II: Analysis of model

# Poissonized model of RHGs: $\mathcal{G}_{\alpha,\nu}(n)$

It is more natural to consider a Poissonized version of  $G_{\alpha,\nu}(n)$ .



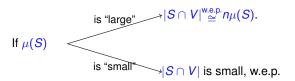
- ▶  $\mathbb{E}|V \cap S|$  is proportional to  $n\mu(S)$  where  $\mu(S) := \iint_{S} f(r,\phi) dr d\phi$ .
- ▶  $|V \cap S_1|$ ,  $|V \cap S_2|$  ... are independent.

Equivalently,  $\forall S \subseteq \mathbb{H}^2$ ,  $|S \cap V| \sim \text{Poisson}(n\mu(S))$ , i.e,  $\forall k > 0$ ,

$$\mathbb{P}(|S\cap V|=k)=e^{-n\mu(S)}\frac{1}{k!}(n\mu(S))^k.$$

4日 > 4周 > 4目 > 4目 > 目 めなの

### Easy (useful) fact



# Easy (useful) fact

$$\text{is } \omega(\underbrace{\frac{\log n}{n}}) \to |S \cap V| \overset{\text{w.e.p.}}{\cong} n\mu(S).$$
 If  $\mu(S)$  otherwise  $|S \cap V| \overset{\text{w.e.p.}}{\leq} (\log n)^{1+o(1)}.$ 

Depoissonization for our purposes (so far) easy

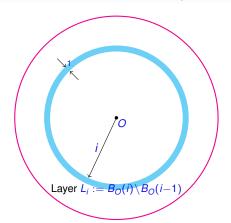
# Henceforth $\frac{1}{2} < \alpha < 1$ .



Do Not Forget!



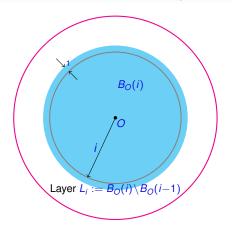
(measure centered balls)



#### Calculations yield[GPP'12]

$$\mu(L_i) \cong \frac{\mu(B_O(i))}{1 - e^{-\alpha}}.$$
 $\mu(B_O(i)) \cong e^{-\alpha(R-i)}.$ 

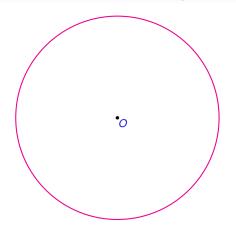
(measure centered balls)



#### Calculations yield[GPP'12]

$$\mu(L_i) \cong rac{\mu(B_O(i))}{1 - e^{-lpha}}.$$
  $\mu(B_O(i)) \cong e^{-lpha(B-i)}.$ 

(measure centered balls)



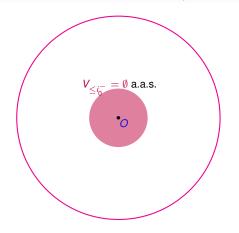
Calculations yield[GPP'12]

$$\mu(L_i) \cong \frac{\mu(B_O(i))}{1 - e^{-\alpha}}.$$
 $\mu(B_O(i)) \cong e^{-\alpha(R-i)}.$ 

Define 
$$V_{\leq i} := V \cap B_O(i)$$
.

Let 
$$i_0 := (1 - \frac{1}{2\alpha})R$$
. So  $\mu(B_O(i_0)) \cong \frac{1}{n}$ .

(measure centered balls)



Calculations yield[GPP'12]

$$\mu(L_i) \cong \frac{\mu(B_O(i))}{1 - e^{-\alpha}}.$$

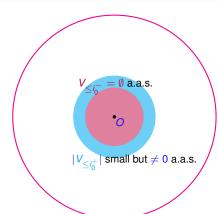
$$\mu(B_O(i)) \cong e^{-\alpha(R-i)}.$$

Define  $V_{\leq i} := V \cap B_O(i)$ .

Let 
$$i_0:=(1-\frac{1}{2\alpha})R$$
. So  $\mu(B_O(i_0))\cong \frac{1}{n}$ .

If 
$$i_0^- = i_0 - \frac{\log R}{\alpha} - \omega(1)$$
, then  $\mathbb{E}|V_{\leq i_0^-}| = n\mu(B_O(i_0^-)) = o(1)$ .

(measure centered balls)



Calculations yield[GPP'12]

$$\mu(L_i) \cong \frac{\mu(B_O(i))}{1 - e^{-\alpha}}.$$

$$\mu(B_O(i)) \cong e^{-\alpha(R-i)}.$$

Define  $V_{\leq i} := V \cap B_O(i)$ .

Let 
$$i_0 := (1 - \frac{1}{2\alpha})R$$
. So  $\mu(B_O(i_0)) \cong \frac{1}{n}$ .

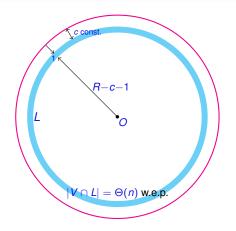
If 
$$i_0^- = i_0 - \frac{\log R}{\alpha} - \omega(1)$$
, then  $\mathbb{E}|V_{\leq i_0^-}| = n\mu(B_O(i_0^-)) = o(1)$ .

If 
$$i_0^+=i_0+\frac{\log R}{\alpha}-\omega(1)$$
, then  $\mathbb{P}(|V_{\leq i_0^+}|>\log n)\leq \frac{1}{\log n}\mathbb{E}|V_{\leq i_0^+}|=o(1)$ .



## Vertices per layer

(measure centered balls)



#### Calculations yield[GPP'12]

$$\mu(L_i) \cong \frac{\mu(B_O(i))}{1 - e^{-\alpha}}.$$

$$\mu(B_O(i)) \cong e^{-\alpha(R-i)}.$$

# Vertex degrees

(measure of non-centered balls)

#### Calculations yield

$$\mu(B_P(R)) = C_{\alpha} e^{-\frac{r_P}{2}} (1 + o(1)).$$

## Vertex degrees

(measure of non-centered balls)

#### Calculations yield

$$\mu(B_P(R)) = C_{\alpha}e^{-\frac{r_P}{2}}(1+o(1)).$$

Thus,

$$\deg(P) = \begin{cases} O(\log n) \text{ (no concentration)}, \\ \text{if } r_P \geq R - 2\log R + O(1), \\ \Theta(ne^{-\frac{r_P}{2}}) \text{ w.e.p.,} \\ \text{otherwise.} \end{cases}$$

### Consequences

- ► A.a.s., a max degree vertex is in  $V_{i_0^+}$  and has degree  $n^{\frac{1}{2\alpha} + o(1)}$  w.e.p.
- ▶ If  $k = C_{\alpha} n e^{-\frac{j}{2}}$ ,  $j \ge i_0^+$ , then w.e.p. the number of degree  $\ge k$  nodes is

$$\cong ne^{-\alpha(R-j)} = n\left(\frac{C_{\alpha}}{k}\right)^{2\alpha}.$$

I.e., power law degree distribution with exponent  $2\alpha + 1$ 

- ► The average degree is  $\pi \nu C_{\alpha}^2 (1 + o(1))$ , i.e., constant!
- ▶ If  $v \notin V_{\leq R-c}$ , c constant,

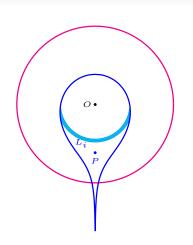
$$\mathbb{P}(\deg(v)=0)\cong C_{\alpha}e^{-c/2}$$

and w.e.p. there are  $\Theta(n)$  such vertices

▶  $V_{\leq R/2}$  induces a clique K (w.e.p.  $|V_{\leq R/2}| = \Theta(n^{1-\alpha})$ )



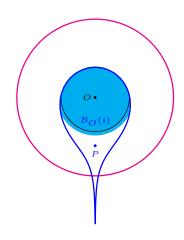
## Location of neighbors of a vertex



#### Calculations yield

$$\mu(\mathcal{B}_{P}(R)\cap L_{i})=\Theta(e^{-\alpha(R-i)}e^{\frac{1}{2}(R-i-r_{P})})$$

## Location of neighbors of a vertex



#### Calculations yield

$$\mu(B_P(R) \cap L_i) = \Theta(e^{-\alpha(R-i)}e^{\frac{1}{2}(R-i-r_P)})$$

As a function of *i* grows like  $e^{(\alpha - \frac{1}{2})i}$ .

#### So, P has:

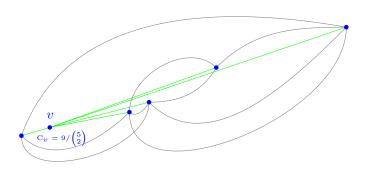
- ▶ more neighbors towards  $\partial B_O(R)$
- ▶ const. fraction of neighbors "near"  $\partial B_O(R)$

◆ロト ◆団 ト ◆ 恵 ト ◆ 恵 ・ 夕 Q ○

#### Visualization of claims

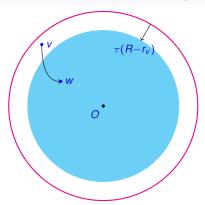
# Non-negligible local clustering coefficient

[GPP'12]



If  $C_{v} := \mathbb{P}_{s,t}(st \in E|s,t \in \mathcal{N}_{v})$ , then  $\mathbb{E}_{v}C_{v} = \Omega(1)$ .

[BFM, EJC'15; FM, AAP'18]

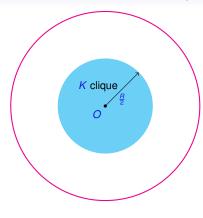


Let  $v \in V$  be s.t.  $R - r_v = \Omega(\log R)$ .

There is a  $\tau > 1$  so that w.e.p.  $\exists w \sim v$  s.t.

$$R-r_w > \tau (R-r_v).$$

[BFM, EJC'15; FM, AAP'18]



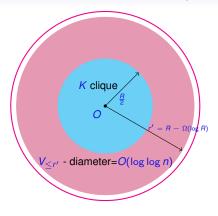
 $|K|_{=}^{\text{wep}}\Theta(n^{1-\alpha})$ 

Let  $v \in V$  be s.t.  $R - r_v = \Omega(\log R)$ .

There is a  $\tau > 1$  so that w.e.p.  $\exists w \sim v$  s.t.

$$R-r_w > \tau(R-r_v).$$

[BFM, EJC'15; FM, AAP'18]



Let  $v \in V$  be s.t.  $R - r_v = \Omega(\log R)$ .

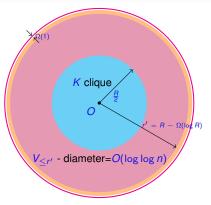
There is a  $\tau > 1$  so that w.e.p.  $\exists w \sim v$  s.t.

$$R-r_w > \tau (R-r_v).$$

$$|K|_{=}^{\text{wep}}\Theta(n^{1-\alpha})$$

$$|V_{\leq r'}|_{=}^{\text{wep}}\Theta(\frac{n}{\text{polylog}(n)})$$

[BFM, EJC'15; FM, AAP'18]



Let 
$$v \in V$$
 be s.t.  $R - r_v = \Omega(\log R)$ .

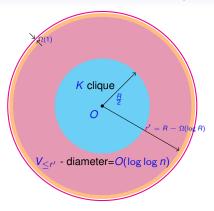
There is a  $\tau > 1$  so that w.e.p.  $\exists w \sim v$  s.t.

$$R-r_w > \tau (R-r_v).$$

$$|K| \stackrel{\mathsf{wep}}{=} \Theta(n^{1-\alpha})$$
$$|V_{\leq r'}| \stackrel{\mathsf{wep}}{=} \Theta(\frac{n}{\text{polylog(n)}})$$

 $|Center component| = \Theta(n)$ 

[BFM, EJC'15; FM, AAP'18]



Let  $v \in V$  be s.t.  $R - r_v = \Omega(\log R)$ .

There is a  $\tau > 1$  so that w.e.p.  $\exists w \sim v$  s.t.

$$R-r_w > \tau (R-r_v).$$

$$|K|_{=}^{\text{wep}}\Theta(n^{1-\alpha})$$

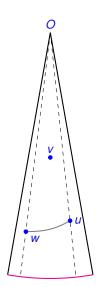
$$|V_{\leq r'}|_{=}^{\text{wep}}\Theta(\frac{n}{\text{polylog(n)}})$$

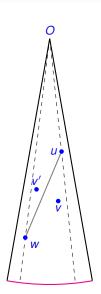
 $|Center component| = \Theta(n)$ 

|2nd component| $\underset{[KM \ 2019]}{\text{wep}}\Theta((\log n)^{1/(1-\alpha)})$ 

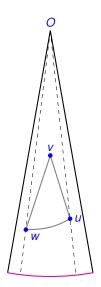


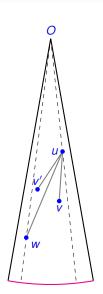
# Forbidden configurations [FK, ICALP'15]



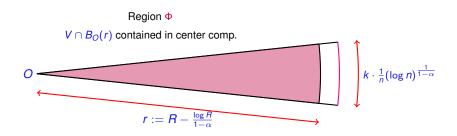


# Forbidden configurations [FK, ICALP'15]





## Idea on upper bound on 2nd component (and diameter)

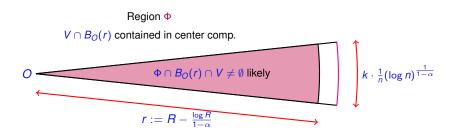


If 
$$k = c(\log n)^{\frac{1}{1-\alpha}}$$
 and  $c$  large enough, then

$$\mathbb{P}(\Phi \cap B_O(R) \cap V = \emptyset) = e^{-\Theta(k)(\log n)^{-\frac{\alpha}{1-\alpha}}} = O(n^{-3}).$$



## Idea on upper bound on 2nd component (and diameter)



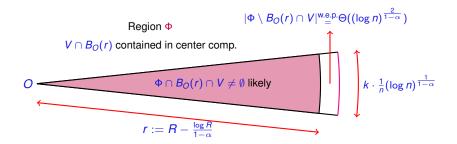
If  $k = c(\log n)^{\frac{1}{1-\alpha}}$  and c large enough, then

$$\mathbb{P}(\Phi \cap B_O(R) \cap V = \emptyset) = e^{-\Theta(k)(\log n)^{-\frac{\alpha}{1-\alpha}}} = O(n^{-3}).$$

A union bound over all  $\Phi$ 's gives  $\mathbb{P}(\exists \Phi, \Phi \cap B_O(R) \cap V = \emptyset) = o(\frac{1}{n})$ .



### Idea on upper bound on 2nd component (and diameter)



If  $k = c(\log n)^{\frac{1}{1-\alpha}}$  and c large enough, then

$$\mathbb{P}(\Phi \cap B_O(R) \cap V = \emptyset) = e^{-\Theta(k)(\log n)^{-\frac{\alpha}{1-\alpha}}} = O(n^{-3}).$$

A union bound over all  $\Phi$ 's gives  $\mathbb{P}(\exists \Phi, \Phi \cap B_O(R) \cap V = \emptyset) = o(\frac{1}{n})$ .

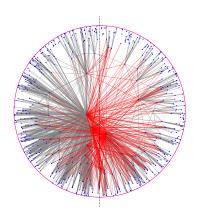


#### Conductance and spectral gap

The graph conductance of the center component H of  $G_{\alpha,\nu}(n)$  is:

$$\varphi(H) := \min_{\substack{S \subseteq V(H) \\ 0 < \operatorname{vol}(S) \le |E(H)|}} \frac{E(S, V(H) \setminus S)}{\operatorname{vol}(S)}.$$

The spectral gap of H is  $\lambda_2(H)$  – the 2nd smallest eigenvalue of the normalized Laplacian of H



By Cheeger's inequality:

$$\frac{1}{2}\varphi^2(H) \le \lambda_2(H) \le 2\varphi(H).$$

Upper bound is almost tight[KM, AAP'18] and

$$\underset{\approx}{\text{wep}} \Theta\left(\frac{1}{n^{2\alpha-1}}\right)$$
 Fairly small!



#### Other ...

- ▶ Bipartite<sup>[KPK, Phys. Rev. E'17]</sup> and higher dimensional analogues, as well as generalizations<sup>[BKL, ESA'17]</sup> have also been considered
- ► Average distance [BKL, arXiv'16]
- ► Separators and treewidth<sup>[BFK, ESA'16]</sup>: Balanced separator hierarchies with separators of size  $O(n^{1-\alpha})$  and  $O(n^{1-\alpha})$  treewidth, a.a.s.
- ▶ Minimum and maximum bisection [KM, AAP'18]
- ► Fast generation<sup>[BKL, ESA'17; vLSMP, ISAAC'15]</sup> and embedding<sup>[BFKL, ESA'16]</sup>
- Connectivity threshold<sup>[BFM, RS&A'16]</sup>
- ▶ Bootstrap percolation<sup>[CF, SP&A'16; KL, ICALP'16; etc.]</sup> in RHGs and GIRGs
- ► Greedy routing [BKLMM, arXiv'17]



#### What next?

(current / near future work)

- How do rumors spread on RHGs?
- How do epidemics/information spread through RHGs? Current work on metastability of contact process
- Work on a dynamic version and establish detection times